

On r -dynamic Coloring of Grids

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Abstract

An r -dynamic k -coloring of a graph G is a proper k -coloring of G such that every vertex in $V(G)$ has neighbors in at least $\min\{d(v), r\}$ different color classes. The r -dynamic chromatic number of a graph G , written $\chi_r(G)$, is the least k such that G has such a coloring. Proving a conjecture of Jahanbekam, Kim, O, and West, we show that the m -by- n grid has no 3-dynamic 4-coloring when $mn \equiv 2 \pmod{4}$. This completes the determination of the r -dynamic chromatic number of the m -by- n grid for all r, m, n .

1 Introduction

When proper graph colorings represent assignment of vertices to categories, in some applications it is desirable for vertices to have neighbors in many categories. This increases the number of colors needed.

An r -dynamic k -coloring is a proper k -coloring f of G such that $|f(N(v))| \geq \min\{r, d(v)\}$ for each vertex v in $V(G)$, where $N(v)$ is the neighborhood of v and $f(U) = \{f(v) : v \in U\}$ for a vertex subset U . The r -dynamic chromatic number, introduced by Montgomery [3] and written as $\chi_r(G)$, is the least k such that G has an r -dynamic k -coloring.

Note that $\chi_1(G) = \chi(G)$, where $\chi(G)$ is the ordinary chromatic number of G . Montgomery called the 2-dynamic chromatic number simply the *dynamic chromatic number*. Many results were motivated by Montgomery's conjecture that $\chi_2(G) \leq \chi(G) + 2$ when G is regular, which remains open. Bounds on χ_r and further references for work on χ_r and its variations appear in [2].

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In this note we complete the solution of a problem in [2]. For $p \in \mathbb{N}$, let $[p] = \{1, \dots, p\}$. The m -by- n grid $G_{m,n}$ is the graph with vertex set $[m] \times [n]$ such that (i, j) and (i', j') are adjacent if and only if $|i - i'| + |j - j'| = 1$. (In more general language, $G_{m,n}$ is the cartesian product of paths with m and n vertices.) The problem of computing $\chi_r(G_{m,n})$ for all r, m, n was proposed in [2].

The following observations are immediate from the definition.

Observation 1.1. $\chi_{r+1}(G) \geq \chi_r(G)$.

Observation 1.2. If $r \geq \Delta(G)$, then $\chi_r(G) = \chi_{\Delta(G)}(G)$.

Observation 1.3. $\chi_r(G) \geq \min\{\Delta(G), r\} + 1$.

To avoid trivialities, assume $m, n \geq 2$. Akbari, Ghanbari, and Jahanbekam [1] proved $\chi_2(G_{m,n}) = 4$. Jahanbekam, Kim, O, and West [2] then determined most of the other values. Since $\Delta(G_{m,n}) \leq 4$, by Observation 1.2 we need only consider $r \leq 4$.

Theorem 1.4 ([2]). *If m and n are at least 2, then*

$$\chi_4(G_{m,n}) = \begin{cases} 4 & \text{if } \min\{m, n\} = 2 \\ 5 & \text{otherwise} \end{cases} \text{ and } \chi_3(G_{m,n}) = \begin{cases} 4 & \text{if } \min\{m, n\} = 2 \\ 4 & \text{if } m \text{ and } n \text{ are both even.} \\ 5 & \text{if } m, n \text{ not both even and } mn \not\equiv 2 \pmod{4} \end{cases}$$

The upper bounds in Theorem 1.4 are by explicit construction. Setting $f(i, j) = i + 2j \pmod{5}$ yields a 4-dynamic 5-coloring of $G_{m,n}$. Constructions for $\chi_r(G_{m,n}) \leq 4$ are obtained by repeating (and truncating when m or n is twice an odd number) the block below.

$$\begin{array}{cccc} 0 & 1 & 2 & 3 \\ 2 & 3 & 0 & 1 \\ 1 & 0 & 3 & 2 \\ 3 & 2 & 1 & 0 \end{array}$$

The lower bounds in Theorem 1.4 are from Observation 1.3, except when m or n is odd and the other is not twice an odd number. In that case, the lower bound follows from the discussion in Lemma 2.1. A statement similar to Lemma 2.1 is used in [2] to prove Theorem 1.4. We include a more explicit version of their observations, because we will use them in proving our main result. The discussion also shows why the remaining case is harder, and it restricts the configurations that need to be considered in that case. For $mn \equiv 2 \pmod{4}$ with $m, n \geq 3$, the authors in [2] proved that five colors suffice and conjectured that five colors are needed. The proof of this conjecture is the content of this note.

Theorem 1.5. *If $m, n \geq 3$ and $mn \equiv 2 \pmod{4}$, then $\chi_3(G_{m,n}) = 5$.*

row is not what we have observed is necessary for the last two columns, so the rows cannot end there. Hence in those congruence classes for n no 3-dynamic 4-coloring exists. \square

Thus we henceforth also assume $n \equiv 2 \pmod{4}$, with $n \geq 6$.

Lemma 2.2. *In the upper left corner, $x_{3,2} = x_{2,3}$, and similarly in the other corners. Furthermore, $m \geq 7$.*

Proof. The cycling of colors as observed in Lemma 2.1 implies the first claim. We then observe that when $m \in \{3, 5\}$, position $(3, 3)$ cannot see three colors. \square

Additional lemmas will restrict the coloring of the interior.

Lemma 2.3. *Every 3-by-2 or 2-by-3 subgrid has vertices of all four colors.*

Proof. By symmetry, it suffices to consider a 2-by-3 subgrid. Suppose that some 2-by-3 subgrid contains at most three colors in a 3-dynamic 4-coloring of $G_{m,n}$. Since the 6-vertex subgraph is bipartite, a color appearing three times would give a vertex three neighbors with the same color, which is forbidden since each vertex has degree at most 4. Hence each of the three colors appears twice. By symmetry, we may assume that it appears as follows.

$$\begin{array}{ccc} a & b & c \\ b & c & a \end{array}$$

Since border vertices have degree at most 3, such a subgrid cannot include a border vertex. Now the fourth color must appear above and below the middle, but it also must appear adjacent to each corner of the rectangle, which puts it on adjacent vertices. \square

We next prove another completely local implication.

Lemma 2.4. *If $x_{i-1,j} = x_{i,j-1}$ and $x_{i,j} \neq x_{i-1,j+1}$, then $x_{i+1,j} = x_{i-1,j+1}$ and $x_{i+1,j+1} = x_{i-1,j}$ (given $i < m$ and $j < n$). The same implication holds with rows and columns exchanged.*

Proof. By symmetry, we may assume $a = x_{i-1,j} = x_{i,j-1}$. Since (i, j) and $(i-1, j+1)$ both see color a and are assumed to have different colors, by symmetry we may let $b = x_{i-1,j+1}$ and $d = x_{i,j}$, as shown below. Now $x_{i,j+1} \notin \{a, b, d\}$, so $x_{i,j+1} = c$, which in turn yields $x_{i+1,j} = b$.

$$\begin{array}{ccc} a & b & \\ a & d & \end{array} \Rightarrow \begin{array}{ccc} a & b & \\ a & d & c \\ & b & \end{array} \Rightarrow \begin{array}{ccc} a & b & \\ a & d & c \\ & b & a \end{array}$$

Now $x_{i+1,j+1} \in \{a, d\}$. If $x_{i+1,j+1} = d$, then having $(i, j+1)$ and $(i+1, j)$ both see a requires $x_{i+2,j}$ and $x_{i,j+2}$ to exist and equal a , but then $(i+1, j+1)$ cannot see a . \square

Lemma 2.5. *Suppose that row r is periodic from columns s to t . If $x_{r+1,s} = x_{r,s+1}$ and $x_{r+1,s+1} = x_{r,s+2}$, then $x_{r+1,j} = x_{r,j+1}$ for $s \leq j < t$ (similarly for columns). This situation cannot occur when $r = 2$ or $s = 2$, or for any r when $x_{r+1,t-1} \neq x_{r,t}$ is known.*

Proof. By symmetry, we may assume that this periodic portion in row r begins d, a, b, c . There is nothing to prove unless $t \geq s + 3$. In that case (shown below), $x_{r,s} \neq x_{r,s+3}$. Lemma 2.3 then implies $x_{r+1,s+2} = x_{r,s+3}$. This establishes the same conditions for the next pair of columns. Continuing to apply Lemma 2.3 copies row r through column t into row $r + 1$, shifted by one column.

$$\begin{array}{c|cccccc} & s & s+1 & & & & t \\ \hline r & d & a & b & c & d & a \\ r+1 & a & b & & & & \end{array} \Rightarrow \begin{array}{c|cccccc} & s & s+1 & & & & t \\ \hline r & d & a & b & c & d & a \\ r+1 & a & b & c & d & a & \end{array}$$

Because $(3, n)$ has only three neighbors, $x_{3,n-1} \neq x_{2,n}$. Also row 2 is periodic up to $t = n$, as shown in Lemma 2.1. Hence the given fact that $x_{r,t}$ does not copy into position $(r + 1, t - 1)$ (in row 2 or later) prevents row $r + 1$ from having two consecutive positions copied from the periodic portion of row r (shifted back by one position). The statement and proof are symmetric for columns. \square

Definition 2.6. A 4-coloring of $G_{m,n}$ is *coherent* if the colors on the 4-by-4 grid in the upper left have the form of the matrix below (under any permutation of the four colors).

$$\begin{array}{cccc} a & b & c & d \\ c & d & a & b \\ b & a & d & c \\ d & c & b & a \end{array}$$

Lemma 2.7. Any 3-dynamic 4-coloring of $G_{m,n}$ is coherent.

Proof. So far we have the pattern of the border positions and their neighbors, as indicated in Lemma 2.1. Applying Lemma 2.5 to row 3 yields $x_{3,3} \neq b$, and applying it to column 3 yields $x_{3,3} \neq c$, so $x_{3,3} = d$. Now Lemma 2.4 yields $x_{4,3} = b$ and $x_{4,4} = a$, and then $x_{3,4} = c$. \square

$$\begin{array}{cccccc} a & b & c & d & a & b \\ c & d & a & b & c & d \\ b & a & d & c & & \\ d & c & b & a & & \\ a & b & & & & \\ c & d & & & & \end{array}$$

3 Proof of Theorem 1.5

We have seen that the colors cycle in the first two rows and in the first two columns. Since the columns have odd length, the third and fourth columns cannot cycle all the way to the bottom. When the first row starts (a, b, c, d) , such cycling would leave the bottom row starting (a, b, c, d) if $m \equiv 1 \pmod{4}$ or (b, a, d, c) if $m \equiv 3 \pmod{4}$. This would contradict the patterns in Lemma 2.1, which showed that these colors are (a, b, d, c) or (b, a, c, d) , respectively. Hence cycling must stop somewhere in column 3 or 4.

Cycling also cannot continue all the way across rows 3 and 4. If it did, then we could delete rows 1 and 2 to obtain a 3-dynamic 4-coloring of $G_{m-2,n}$, contradicting the minimality of m and thereby completing the proof.

We introduce definitions to facilitate study of where cycling of colors breaks down.

Definition 3.1. Given $(i, j) \in [m] \times [n]$, define $(p, q) \in [4] \times [4]$ by $i \equiv p \pmod{4}$ and $j \equiv q \pmod{4}$. A position (i, j) is *correct* if $x_{i,j} = x_{p,q}$. For $i > 4$, positions (i, j) and $(i + 1, j)$ are *flipped* if they are incorrect but $x_{i,j} = x_{p+1,q}$ and $x_{i+1,j} = x_{p,q}$. Similarly, for $j > 4$, positions (i, j) and $(i, j + 1)$ are *flipped* if $x_{i,j} = x_{p,q+1}$ and $x_{i,j+1} = x_{p,q}$.

We observed at the beginning of this section that neither rows 3 and 4 nor columns 3 and 4 are completely correct. Hence there is a first column having an incorrect position in row 3 or 4, and there is a first row having an incorrect position in column 3 or 4. We next show that these first incorrect positions are in row 3 and column 3, respectively. Note that correctness requires $x_{3,j} = x_{2,j+1}$ when j is even, but these positions have different colors when j is odd. We also show that the first violation of correctness leads to flipped positions.

Lemma 3.2. *Given $4 \leq 2s \leq n - 2$, suppose that row $2r$ is periodic through column t such that $t > 2s$ and $x_{2r+1,t-1} \neq x_{2r,t}$. If rows $2r + 1$ and $2r + 2$ are correct through column $2s$, then either they are both correct through column $2s + 2$, or positions $(2r + 2, 2s + 1)$ and $(2r + 2, 2s + 2)$ are correct while positions $(2r + 1, 2s + 1)$ and $(2r + 1, 2s + 2)$ are flipped. The same statement holds with the roles of rows and columns switched.*

Proof. Let the cycle of colors in row $2r$ have (a, b, c, d) ending at column $2s$.

	$2s$					
$2r$	a	b	c	d	a	b
	c	d	a	b	c	d
	b	a	d	c	b	a
	d	c	b	a	d	c

Since $(2r + 1, 2s)$ is correct and $t > 2s$, we have $x_{2r+1,2s} = x_{2r,2s+1} = a$ and $x_{2r,2s+2} = b$. By Lemma 2.5, $x_{2r+1,2s+1} \neq b$. Whether $(2r + 1, 2s + 1)$ is correct or not, Lemma 2.4 now yields $x_{2r+2,2s+1} = b$ and $x_{2r+2,2s+2} = a$, which makes $(2r + 2, 2s + 1)$ and $(2r + 2, 2s + 2)$ correct.

If $x_{2r+1,2s+1} = d$, then $(2r + 1, 2s + 2)$ sees $\{a, b, d\}$. Hence $x_{2r+1,2s+2} = c$, making $(2r + 1, 2s + 1)$ and $(2r + 1, 2s + 2)$ also correct.

Otherwise, $x_{2r+1,2s+1} = c$. Now $(2r + 1, 2s + 1)$ sees $\{a, b, c\}$, so $x_{2r+1,2s+1} = d$ and positions $(2r + 1, 2s + 1)$ and $(2r + 1, 2s + 2)$ are flipped. □

Applying Lemma 3.2 with $r = 1$ yields a value s (with $s \geq 2$) such that row 3 is correct through column $2s$ with $(3, 2s + 1)$ and $(3, 2s + 2)$ flipped, and row 4 is correct through column $2s + 2$. Applying the column version with $s = 1$ also yields a value r (with $r \geq 2$) such that column 3 is correct through row $2r$ with $(2r + 1, 3)$ and $(2r + 2, 3)$ flipped, and column 4 is correct through row $2r + 2$.

Definition 3.3. With respect to a coherent block in the upper left, an (r, s) -*partial coloring* of a grid is a vertex 4-coloring that flips positions $(3, 2s+1)$ and $(3, 2s+2)$ and flips positions $(2r+1, 3)$ and $(2r+2, 3)$, but is correct on all of the following: the first three rows through column $2s$, row 4 through column $2s+2$, the first three columns through row $2r$, and column 4 through row $2r+2$.

We remarked before Definition 3.3 that every 3-dynamic 4-coloring of $G_{m,n}$ is an (r, s) -partial coloring for some (r, s) . Hence our next lemma completes the proof of the theorem.

Lemma 3.4. For $r, s \geq 2$, an (r, s) -*partial coloring* of a grid cannot be completed to a 3-dynamic 4-coloring.

Proof. We use induction on $r + s$. For $r = s = 2$, we have the coloring shown below. It cannot be completed, because $(5, 4)$ requires $x_{5,5} = c$, but $(4, 5)$ requires $x_{5,5} = b$.

$$\begin{array}{cccccc}
 a & b & c & d & & \\
 c & d & a & b & & \\
 b & a & d & c & a & b \\
 d & c & b & a & d & c \\
 & & a & d & & \\
 & & c & b & &
 \end{array}$$

Now suppose $r + s > 4$. Since $(3, 2s+1)$ and $(3, 2s+2)$ are flipped, $x_{4,2s} = x_{3,2s+1}$. Since $(4, 2s+1)$ and $(4, 2s+2)$ are correct, $x_{4,2s+1} \neq x_{3,2s+2}$. Hence Lemma 2.4 yields $x_{5,2s+1} = x_{3,2s+2}$ and $x_{5,2s+2} = x_{3,2s+1}$, which means that $(5, 2s+1)$ and $(5, 2s+2)$ are flipped. Similarly, $(2r+1, 5)$ and $(2r+2, 5)$ are flipped.

	1	2	3	2s								
1												
2												
3			a	b	c	d	a	b	c	d	b	a
4			c	d	a	b	c	d	a	b	c	d
5			b	a						a	b	
			d	c								
			a	b								
2r			c	d								
			d	a	b							
			b	c	d							

In the matrix above, we have applied a permutation of the colors to illustrate the smaller instance of the problem to which we will apply the induction hypothesis. Assume these labels. If $x_{5,5} = b$, then Lemma 2.5 applies in row 5 to contradict $x_{5,2s+1} \neq x_{4,2s+2}$. If $x_{5,5} = c$, then Lemma 2.5 applies in column 5 to contradict $x_{2r+1,5} \neq x_{2r+2,4}$. Hence $x_{5,5} = d$. Now Lemma 2.4 yields $x_{6,5} = b$, $x_{6,6} = a$, and $x_{5,6} = c$. Furthermore, the positions in $\{5, 6\} \times \{5, 6\}$ are correct. (This argument also implies that $r > 2$ if and only if $s > 2$.)

We now build a smaller instance of the problem. Since positions $(5, 2s+1)$ and $(5, 2s+2)$ are flipped and $x_{5,2s+1} \neq x_{4,2s+2}$, Lemma 3.2 implies that for some s' with $s' \leq s$, rows 5 and

6 are correct through column $2s'$, with positions $(6, 2s' + 1)$ and $(6, 2s' + 2)$ correct while $(5, 2s' + 1)$ and $(5, 2s' + 2)$ are flipped. The same discussion applies to columns 5 and 6, yielding r' with $r' \leq r$. Note that $r', s' \geq 2$.

Let M be the matrix obtained by deleting the first two rows and first two columns of the given coloring. Since the positions in $\{5, 6\} \times \{5, 6\}$ are correct, M is coherent. Hence M is an $(r' - 2, s' - 2)$ -partial coloring of a grid. By the induction hypothesis, it cannot be extended to a 3-dynamic 4-coloring. Since M extends to row m and column n of the original coloring, we conclude that the original coloring also does not extend to a 3-dynamic 4-coloring. \square

References

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